

InfERL

Scalable and extensible static analysis for Erlang

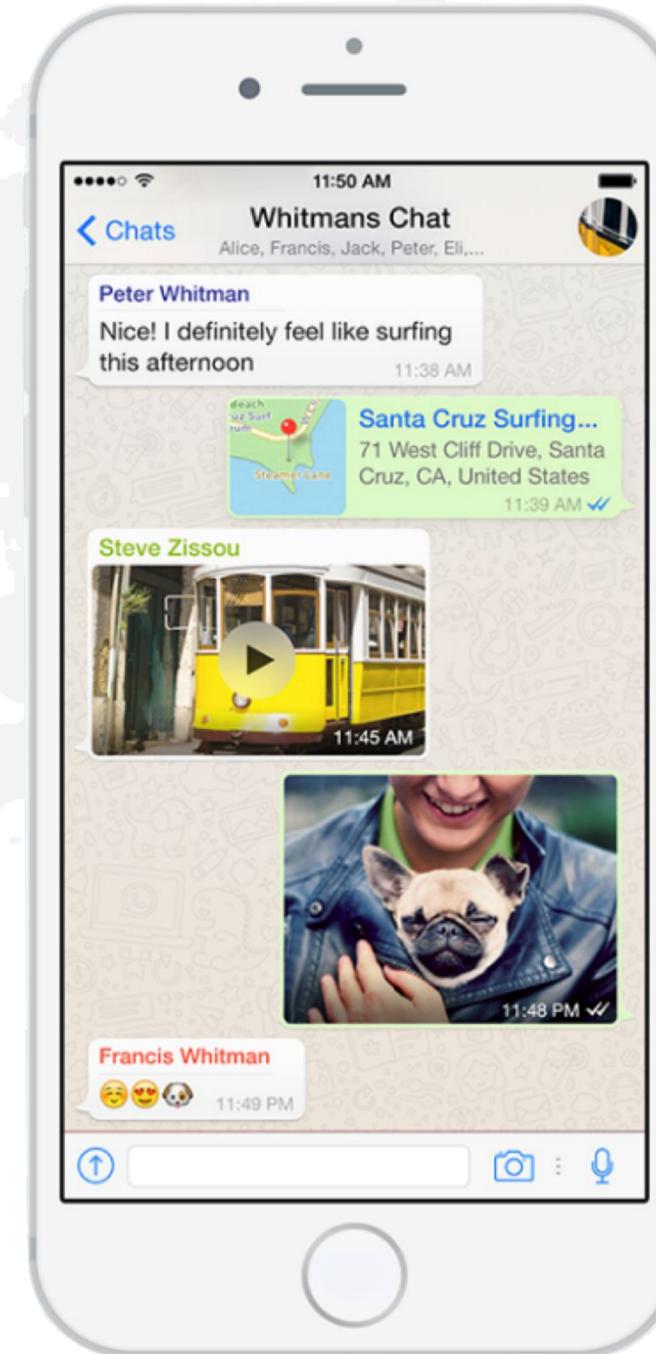
Ákos Hajdu, Matteo Marescotti, Thibault Suzanne,
Ke Mao, Radu Grigore, Per Gustafsson, Dino Distefano
WhatsApp Dev Infra





2 Billion

People around the world use WhatsApp daily





100 Billion

Messages daily

Simple

So anyone can use it.

Reliable

So that messages go through no matter what.

End-to-end Encrypted

So only sender and receiver of the message can see its content.

WhatsApp server

- Millions of lines of **Erlang** code
- **Static analysis** applied so far
 - Linters
 - eqWAlizer (type checker)
 - github.com/WhatsApp/eqwalizer
 - (Incremental) Dialyzer
 - github.com/erlang/otp/pull/5997

WhatsApp/ eqwalizer

A type-checker for Erlang

8 Contributors 4 Issues 383 Stars 10 Forks

github.com
GitHub - WhatsApp/eqwalizer: A type-checker for Erlang
A type-checker for Erlang. Contribute to WhatsApp/eqwalizer development by creating an account on GitHub.



erlang/otp

#5997 dialyzer: Add incremental analysis mode



12 comments 11 reviews 44 files +4911 -882

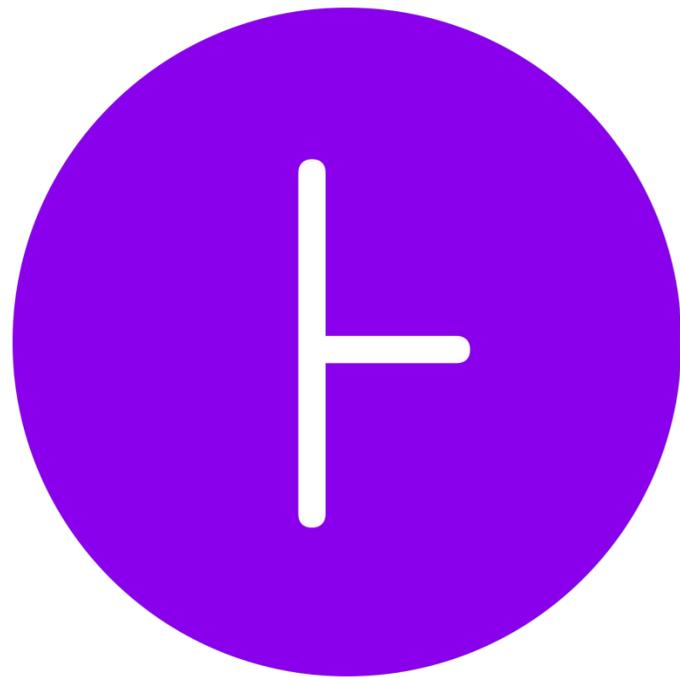
TD5 • May 18, 2022 1 commit

github.com

dialyzer: Add incremental analysis mode by TD5 · Pull Request #5997 ...
Implements a new dialyzer --incremental mode. Incremental mode primarily differs from the previous, "classic", ways of running Dialyzer, ...



Infer



Open-source static analysis platform

fbinfer.com

github.com/facebook/infer

Developed at Meta

Runs on tens of thousands of code changes monthly, reporting thousands of issues

Language frontends

C, C++, Objective-C, Java, C#, **Erlang**

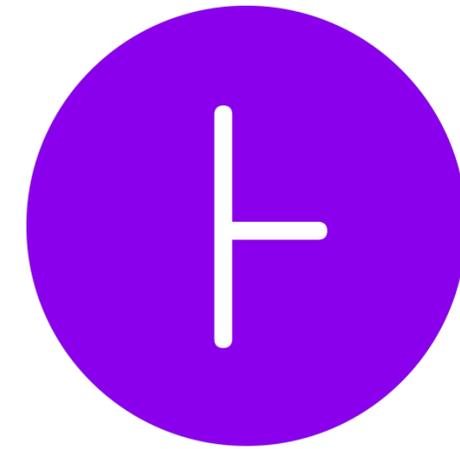
20+ analyses

Memory safety, data races, time complexity, deadlocks, temporal properties, ...

This talk



+



01 Examples

02 Challenges

03 Compilation and Analysis

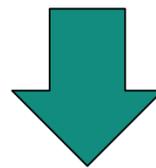
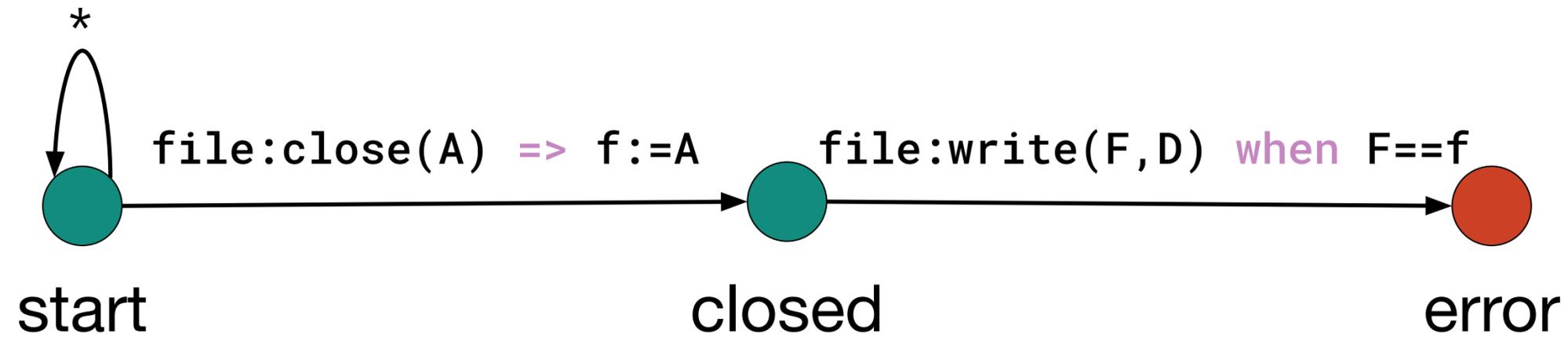
04 Results



Examples

User-specified properties

Example: file:write after file:close



writep.top1

```
1 property WriteAfterClose
2   start -> start: *
3   start -> closed: "file:close/1"(A,Ret) => f:=A
4   closed -> error: "file:write/2"(F,D,Ret) when F==f
```

Example: file:write after file:close

Interprocedural

write.erl

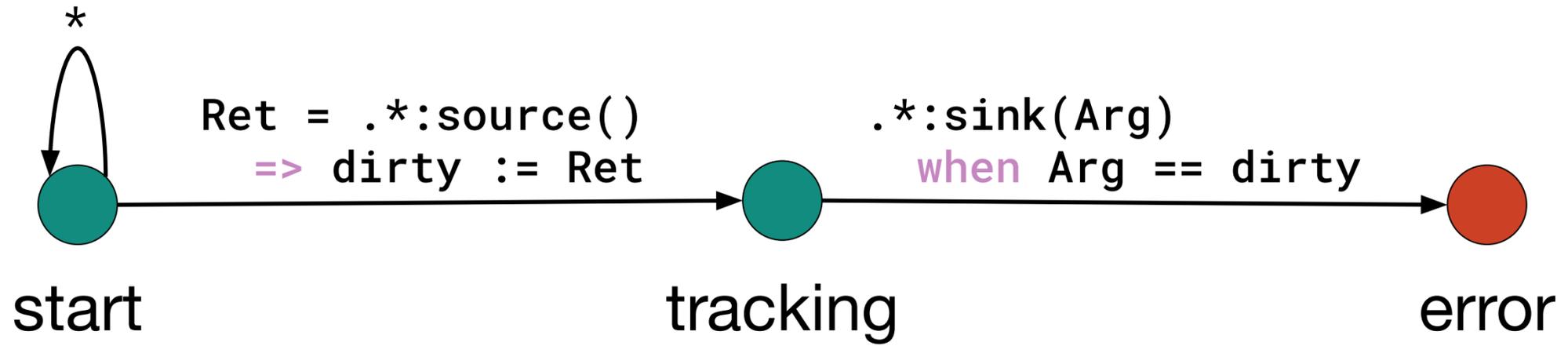
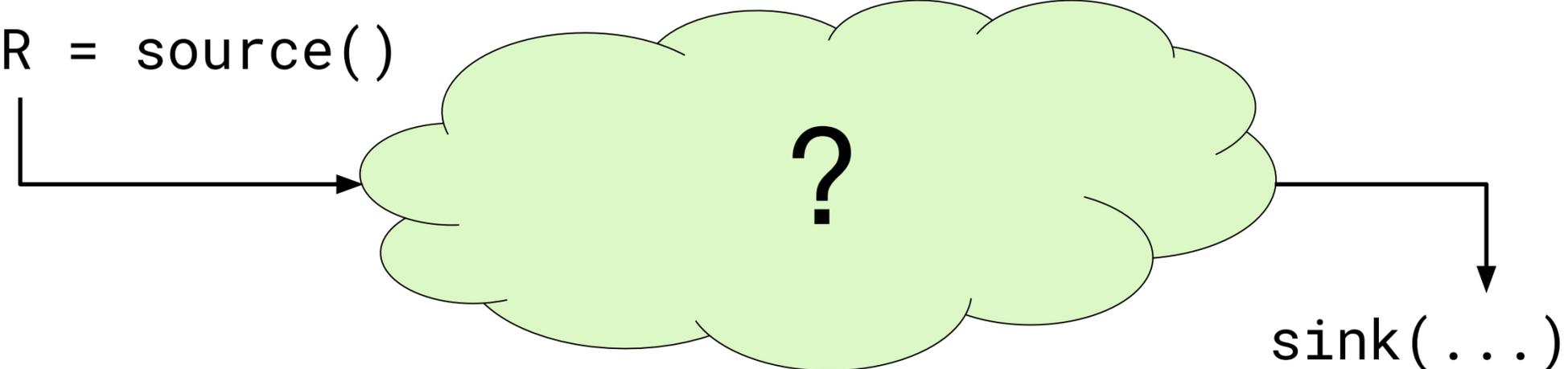
```
1 -module(write).
2
3 good(F) -> nop(F), file:write(F, "hi").
4 bad(F)   -> op(F),  file:write(F, "hi").
5 nop(_)   -> ok.
6 op(F)    -> file:close(F).
```

```
$ infer --top1-only --top1-properties writep.top1 -- erlc write.erl
```

```
write.erl:4: error: Top1 Error
property WriteAfterClose reaches state error.
```

```
3. good(F) -> nop(F), file:write(F, "hi").
4. bad(F)   -> op(F),  file:write(F, "hi").
   ^
```

Example: taint analysis

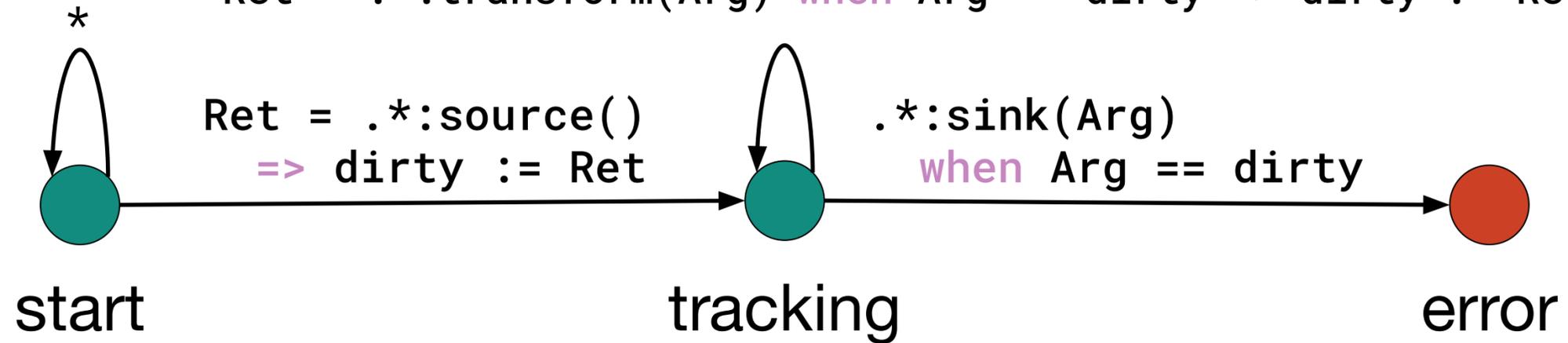


Example: taint analysis with transformations

Built-in data structures, e.g:
{X, Y} is `__erlang_make_tuple(X, Y)`
[H | T] is `__erlang_make_cons(H, T)`

Add your functions of interest

```
Ret = __erlang_make_.*(A)           when A == dirty => dirty := Ret
Ret = __erlang_make_.*(A, B)        when A == dirty => dirty := Ret
Ret = __erlang_make_.*(A, B)        when B == dirty => dirty := Ret
Ret = __erlang_make_.*(A, B, C)     when A == dirty => dirty := Ret
Ret = __erlang_make_.*(A, B, C)     when B == dirty => dirty := Ret
Ret = __erlang_make_.*(A, B, C)     when C == dirty => dirty := Ret
*
Ret = .*:transform(Arg) when Arg == dirty => dirty := Ret
```





Examples

Reliability issues

Example - Reliability

sample.erl

```
1 -module(sample).
2
3 head([]) -> empty;
4 head([H|_]) -> {ok, H}.
5
6 good() ->
7   {ok, 1} = head([1, 2]).
8
9 bad() ->
10  empty = head([1, 2]).
11
12 bad2() ->
13  {ok, 3} = head([1, 2]).
14
15 very_bad() ->
16  head(123).
```

```
$ infer --pulse -- erlc sample.erl
```

```
sample.erl:3: error: No Matching Function Clause
no matching function clause at line 3, column 1.
```

```
sample.erl:16:3: calling context starts here
```

```
15. very_bad() ->
16.   head(123).
    ^
```

```
sample.erl:16:3: in call to `head/1`
```

```
15. very_bad() ->
16.   head(123).
    ^
```

```
sample.erl:3:1: no matching function clause here
```

```
3. head([]) -> empty;
   ^
```

```
sample.erl:10: error: No Match Of Rhs
no match of RHS at line 10, column 3.
```

```
9. bad() ->
10.   empty = head([1, 2]).
    ^
```

```
sample.erl:13: error: No Match Of Rhs
no match of RHS at line 13, column 3.
```

```
12. bad2() ->
13.   {ok, 3} = head([1, 2]).
    ^
```

Reliability issues

```
1 % Bad key
2 M = #{},
3 M#{2 := 3}.
4
5
6 % Bad map
7 L = [1, 2, 3],
8 L#{1 => 2}.
9
10
11 % Bad record
12 R = #rabbit{name="Pogi"},
13 R#person.name.
14
15
16 % No matching branch in try
17 tail(X) ->
181 try X of
9     [_ | T] -> {ok,T}
20 catch
21     _ -> error
22 end.
23 % No matching case clause
24 tail(X) ->
25     case X of
26         [_ | T] -> T
27     end.
28
29
30 % No matching function clause
31 tail([_ | Xs]) -> Xs.
32
33
34 % No match of rhs
35 [H | T] = [].
36
37
38 % No true branch in if
39 sign(X) ->
40     if
41         X > 0 -> pos;
41         X < 0 -> neg
43     end.
44
```



Challenges



Context

Support for Java, C/C++/Obj-C, C#, ...

Built-in analyzers: Pulse, TopI, ...

Compositional, interprocedural

Higher-order functions, closures

Recursion

Captured variables, scopes

Functional

```
1 map(_, []) -> [];  
2 map(F, [H | T]) -> [F(H) | map(F, T)].  
3  
4 main() ->  
5     X = 1,  
6     map(fun(Y) -> Y + X end, [1, 2, 3]).
```

Let it crash

User vs runtime exceptions

Fault tolerance with supervisors

```
1 head([H | _]) -> H.
```

Dynamic typing

Dynamic types used pervasively

Pattern matching seldom complete

```
1 len([]) -> 0;  
2 len([_ | T]) -> 1 + len(T).
```

Concurrency

Designed with concurrency in mind

Primitives for send/receive

High-level abstractions

```
1 receive
2   {From, X, Y} -> From ! X + Y
3   {From, X}   -> From ! X + 1
4 end
```

Scalability

WhatsApp server is huge

Early signals for developers

Erlang-specific abstractions

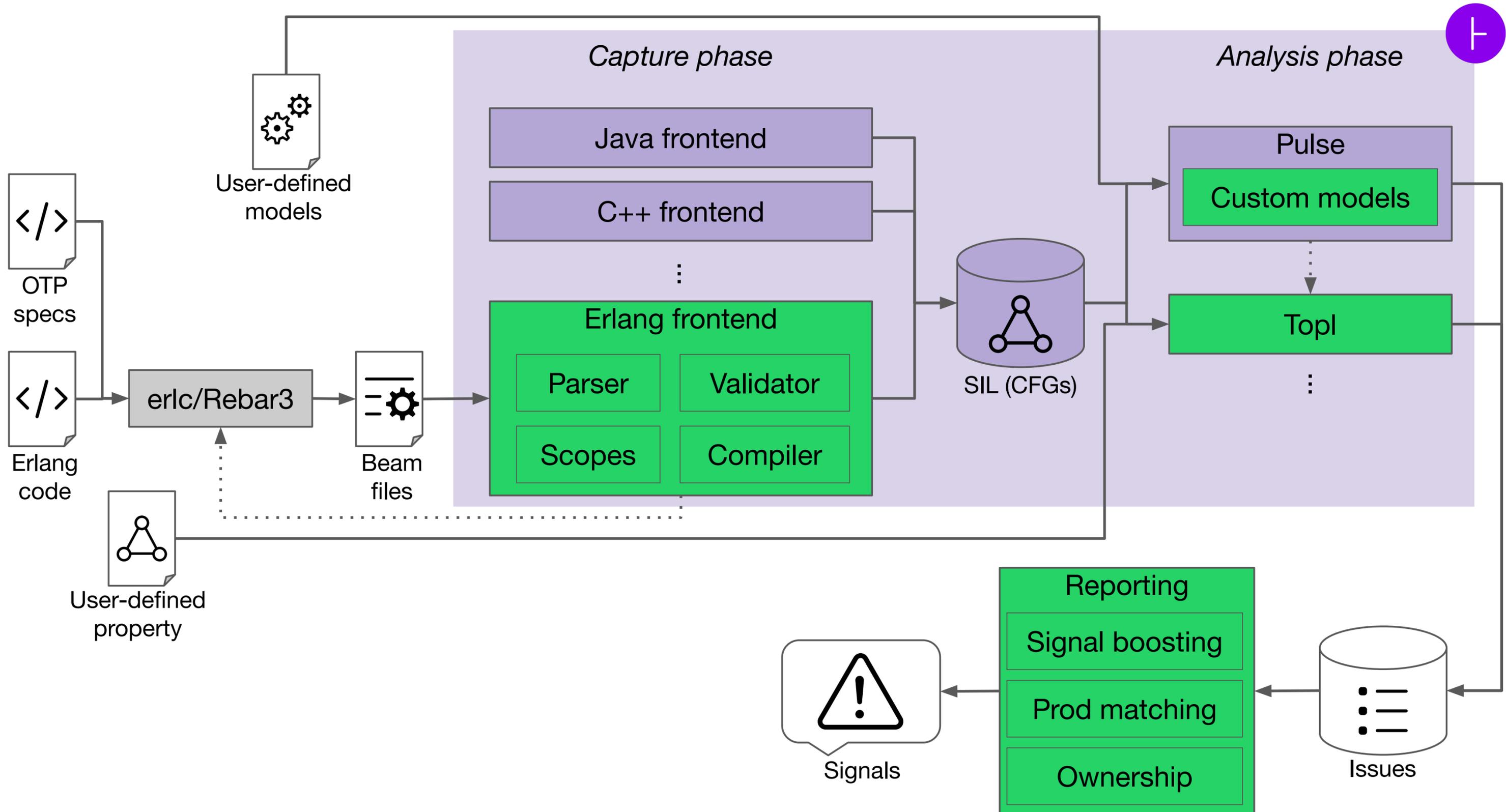
```
1 [X * Y || X <- [1, 2, 3],  
2   Y <- [1, 2, 3]]
```



Compilation and Analysis

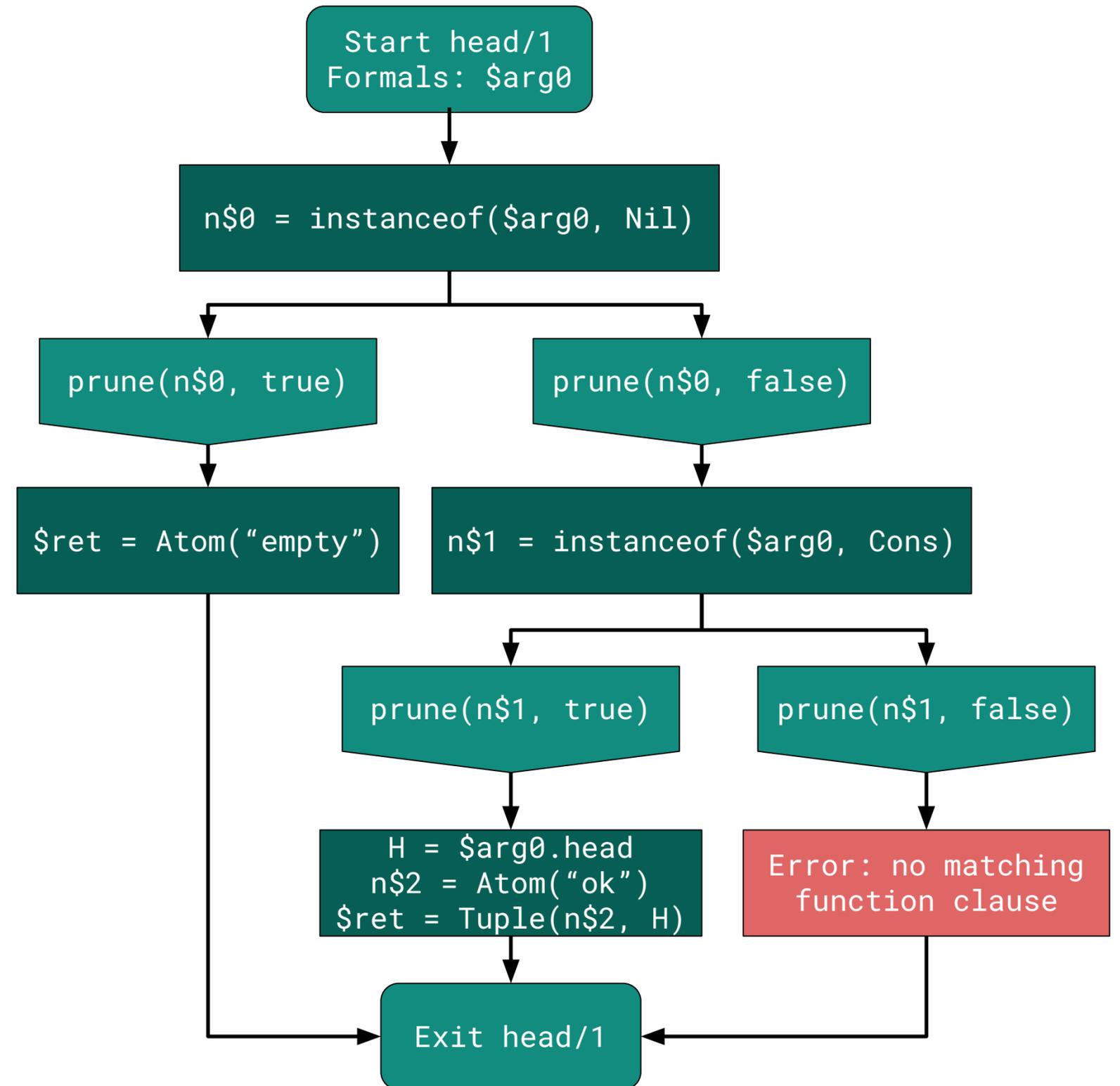
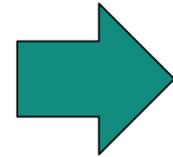
Basics

Overview



Compilation basics

```
1 head([]) -> empty;  
2 head([H|_]) -> {ok, H}.
```



Analysis basics

- Summaries

- Compute once per function
- Compact representation of interesting behaviors

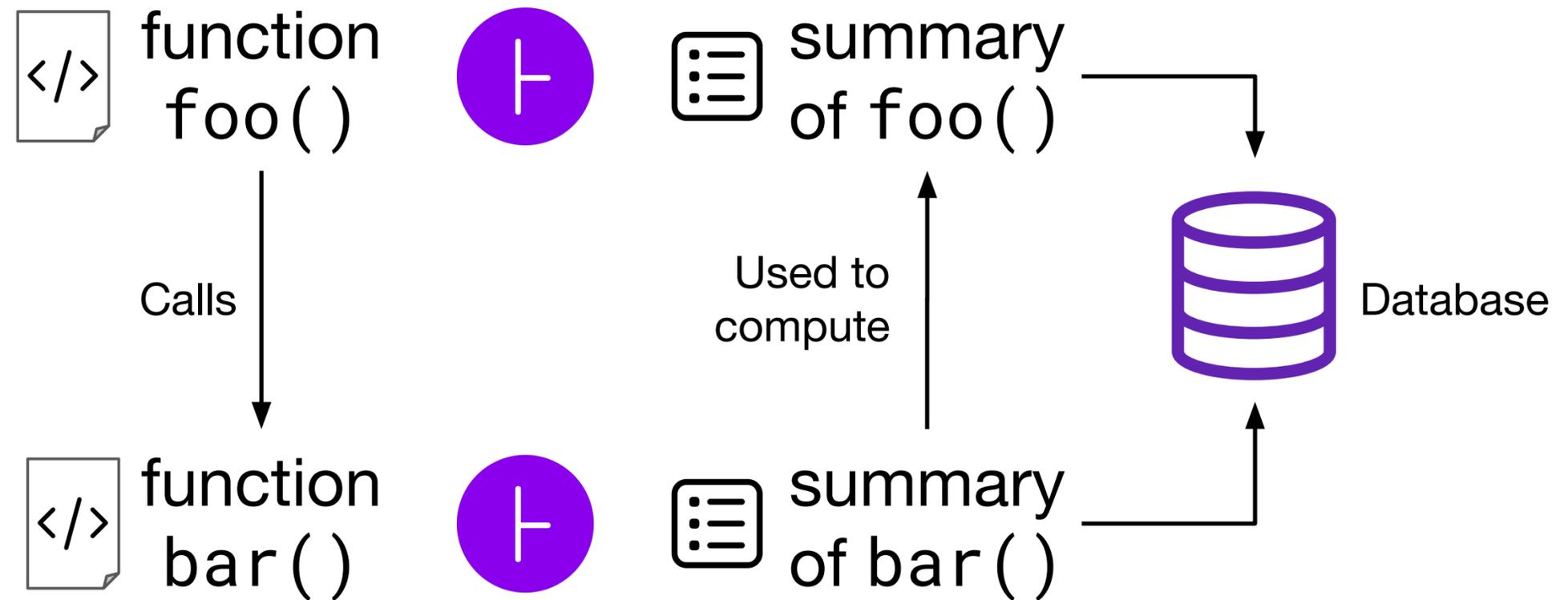
```
1 head([]) -> empty;  
2 head([H|_]) -> {ok, H}.
```



1. **pre:** \$arg0 instanceof Nil
post: \$ret = Atom("empty")
2. **pre:** \$arg0 instanceof Cons
post: \$ret = Tuple(Atom("ok"), \$arg0.head)
3. **pre:** not \$arg0 instanceof Nil and
not \$arg0 instanceof Cons
post: ERROR

Analysis basics

- **Modular:** analyze one procedure (+deps) at a time
- **Compositional:** summary can be used in all calling contexts



Analysis basics

- Under the hood: **Pulse** analyzer (+TopI)
 - Incorrectness separation logic
 - Under-approximate

[Pre] Code [ok:Res]
[Pre] Code [err:Res]
[[Res]] \subseteq
[[Code]]_{ok/err} ([[Pre]])

Incorrectness Logic

PETER W. O'HEARN, Facebook and University College London, UK

Program correctness and incorrectness are two sides of the same coin. As a programmer like to have correctness, you might find yourself spending most of your time reasoning about incorrectness. This includes informal reasoning that people do while looking at or thinking about their code, supported by automated testing and static analysis tools. This paper describes a separation logic for reasoning about incorrectness which is, in a sense, dual to the logic of correctness.

CCS Concepts: • Theory of computation

Additional Key Words and Phrases

ACM Reference Format:

Peter W. O'Hearn. 2020. Incorrectness Logic. *ACM SIGPLAN Notices*, 55(12), 32 pages. <https://doi.org/10.1145/3621451>

1 INTRODUCTION

When reasoning in formal logic

Finding Real Bugs in Big Programs with Incorrectness Logic

QUANG LOC LE, University College London and Meta, UK

AZALEA RAAD, Imperial College London and Meta, UK

JULES VILLARD, Meta, UK

JOSH BERDINE, Meta, UK

DEREK DREYER, MPI-SWS, Germany

PETER W. O'HEARN, Meta and University College London, UK

Local Reasoning about the Presence of Bugs: Incorrectness Separation Logic

Azalea Raad¹, Josh Berdine², Hoang-Hai Dang¹,
Derek Dreyer¹, Peter O'Hearn^{2,3}, and Jules Villard²

¹University of Applied Sciences (MPI-SWS),
Mannheim, Germany
²Imperial College London, UK
³University College London, UK

of work on local reasoning for
or proving their *presence*. We
reasoning about the presence of
relations: 1) separation logic and
theory of this new *incorrectness*



Compilation and Analysis

Details

Type specs

- Avoid false positives due to potential non-exhaustive pattern matching

```
1 -spec len(list()) -> integer().
2 len([]) -> 0;
3 len([_ | T]) -> 1 + len(T).
```

- Return value of unknown functions (from OTP)

```
1 min(Xs) ->
2   case lists:sort(Xs) of
3     [] -> error;
4     [X|_] -> {ok, X}
5   end.
- spec sort(list()) -> list().
```

Data structures

- Tuples, records: **full support**
- Lists: **approximate**, loop unrolling bound
- Maps: recency **abstraction**

```
1 -spec ok(map()) -> any().
2 ok(M) ->
3   case maps:is_key(key, M) of
4     true -> maps:get(key, M);
5     _ -> nope
6   end.
```



```
1 -spec bad(map()) -> any().
2 bad(M) ->
3   maps:get(key, M).
4
5
6
```



Closures, higher-order functions, recursion

- Currently

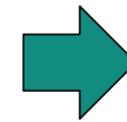
- Limited support
- Scope analysis: captured vs local
 - github.com/erlang/otp/issues/5379

```
1 (X=1)
2 +
3 (begin F=(fun () -> X=2 end), F() end)
```

- Plans

- Unknown closures
 - Specialization*
 - Defunctionalization**
- Recursion
 - Fixed-point computation
 - Custom models (for standard library)

```
1 g(F) -> F().
2
3 f() -> 1.
4
5 main() -> g(fun f/0).
```



```
1 g_f() -> f().
2
3
4
5 main_f() -> g_f().
```

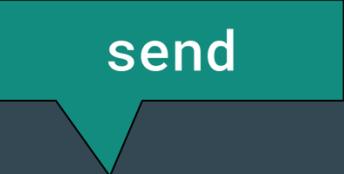
*Corentin De Souza: Higher-order function specialization in Infer, talk at Infer workshop 2022.

**John C Reynolds. 1972. Definitional interpreters for higher-order programming languages. Proc. of the ACM annual conf. Vol. 2. 717–740. 31

Concurrency - send/receive

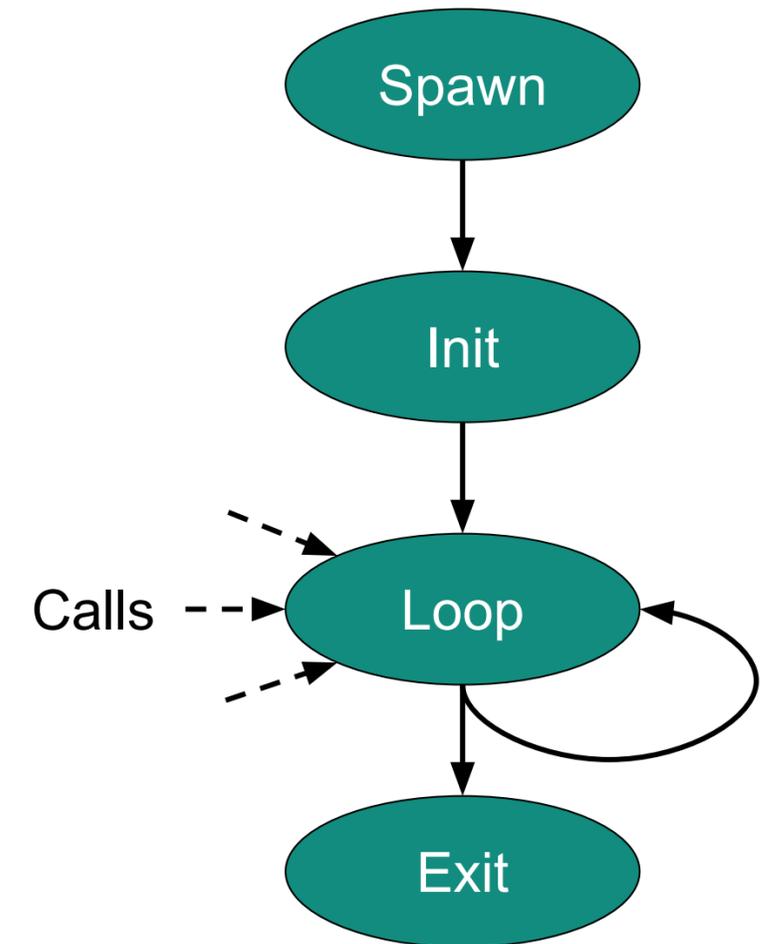
- Currently
 - send: no-op
 - receive: nondeterministic
- Plans
 - Generalization of function call
 - Connect send/receive
 - Create summaries
 - Fixed-point computation

```
1 -module(concurrent).  
2 -export([f/0]).  
3  
4 f() ->  
5   receive  
6     {From, X} when is_integer(X) -> From ! X + 1;  
7     {From, _} -> From ! oops  
8   end,  
9   f().
```



Concurrency - high-level abstractions

- Not yet supported
- Example plan: `gen_server`
 - Distributed objects/actors
 - Treat as objects (Infer has support)



Library functions

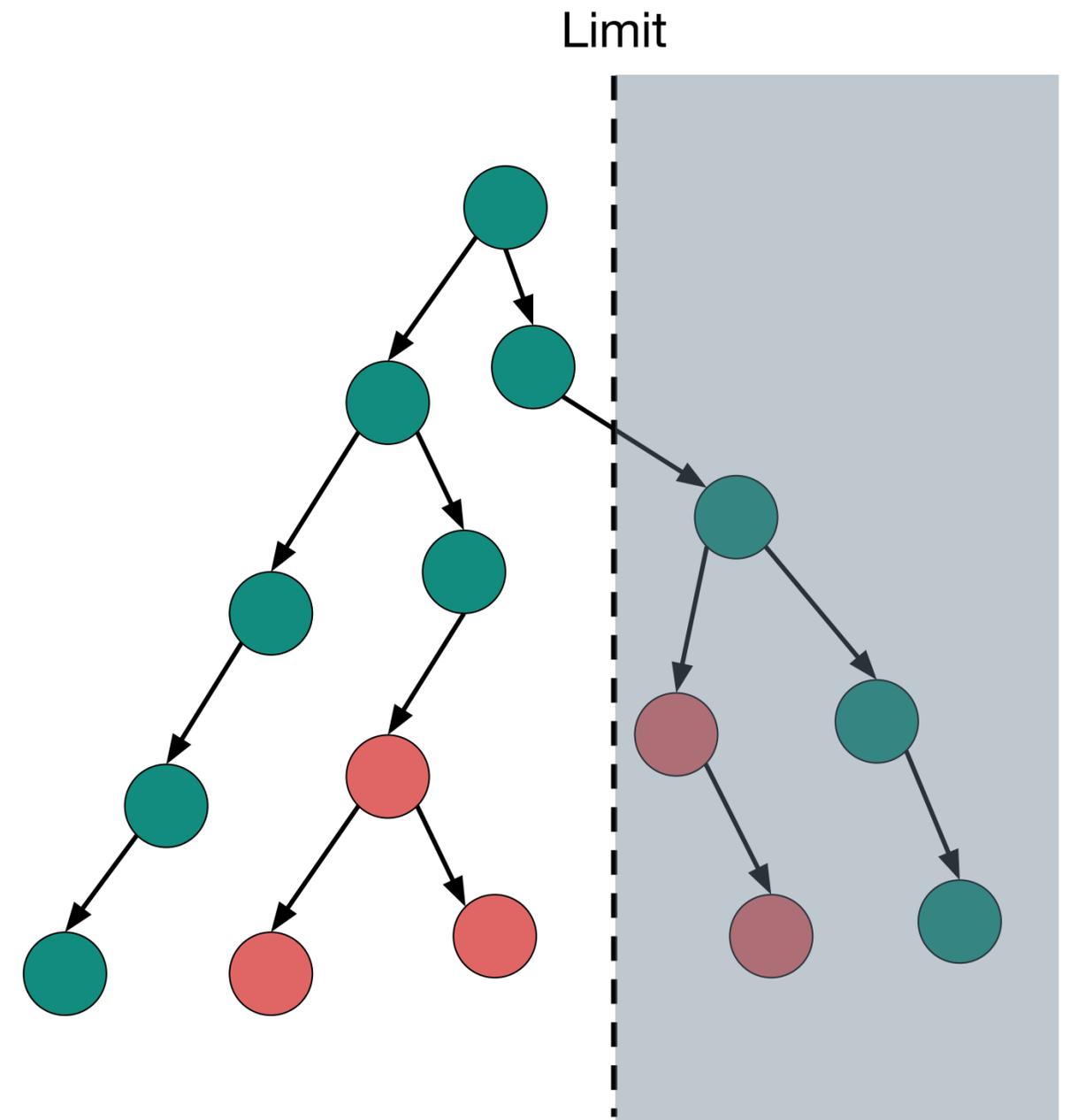
- OTP: use **type specs**
- **User-defined** models
 - Plan: more behaviors
 - Plan: more user-friendly

```
1 {
2   "pulse-models-for-erlang": [
3     {
4       "selector": [
5         "MFA",
6         {
7           "module": "some_module",
8           "function": "complicated_function",
9           "arity": 0
10        }
11      ],
12      "behavior": [
13        "ReturnValue",
14        [
15          "Tuple",
16          [
17            [ "Atom", "true" ],
18            null
19          ]
20        ]
21      ]
22    }
23  ]
24 }
```

```
-module(some_module).
complicated_function() ->
{true, nondet()}.
```

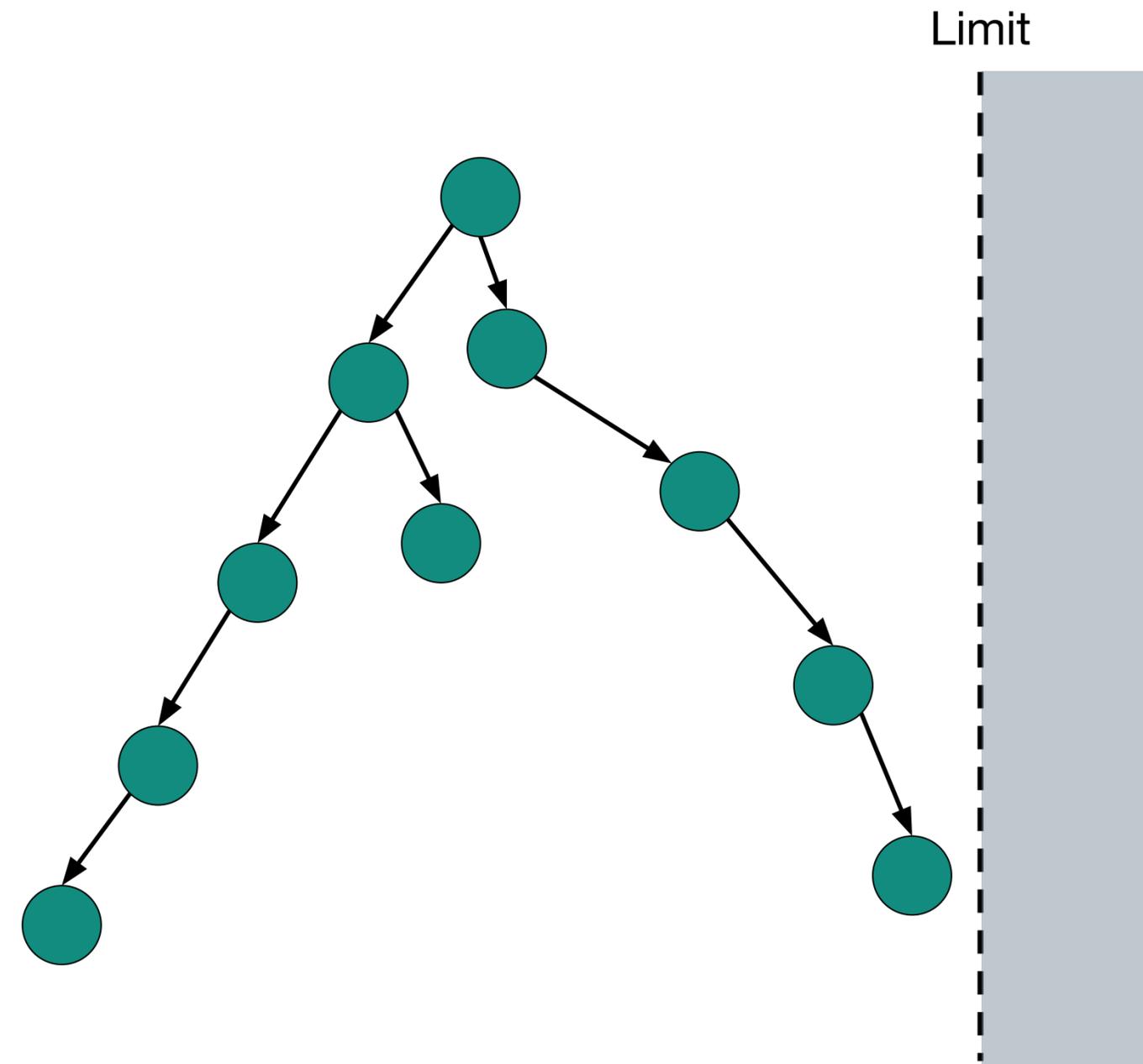
TopI

- Analyzer for **user-defined** (temporal) **properties**
 - Extends Pulse
- Scalability **issues**
 - Hit internal limit during exploration
 - False negatives
- **Improvements**
 - More expensive normalization
 - Triggered selectively
 - Dynamic types in solver
 - Garbage collector



TopI

- Analyzer for **user-defined** (temporal) **properties**
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Results

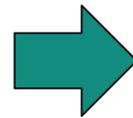
Test functions

- ~750 small functions
 - github.com/facebook/infer/tree/main/infer/tests/codetoanalyze/erlang

```
1 test_add_Ok() ->
2   X = 2,
3   Y = 3,
4   case X + Y of
5     5 -> ok;
6     _ -> ?CRASH
7   end.
```



```
1 test_add_Bad() ->
2   X = 2,
3   Y = 3,
4   case X + Y of
5     5 -> ?CRASH;
6     _ -> ok
7   end.
```



Compile & run	Infer	
Ok	Ok	True negative
Ok	Issue	False positive
Crash	Ok	False negative
Crash	Issue	True positive

WhatsApp server

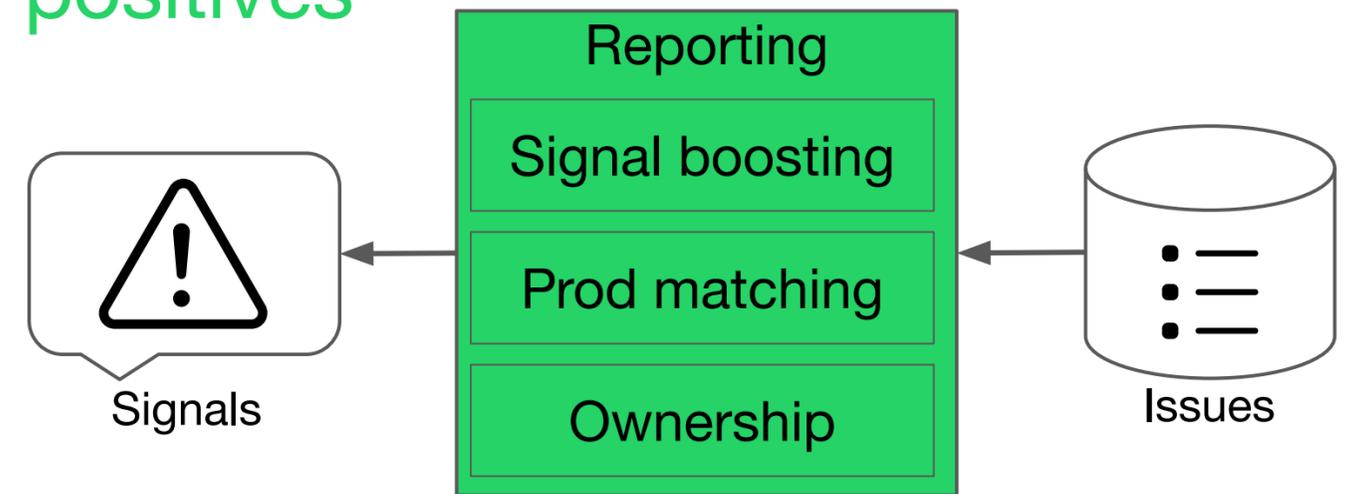
- Initial deployment with **reliability: many false positives**

- **Taint** properties + prod matching

- 200 found
- 21 surfaced
- Fixed: **2 high-pri, 1 very high-pri**

- Signal boosting

- **Dynamic** analysis: FAUSTA



2022 IEEE Conference on Software Testing, Verification and Validation (ICST)

FAUSTA: Scaling Dynamic Analysis with Traffic Generation at WhatsApp

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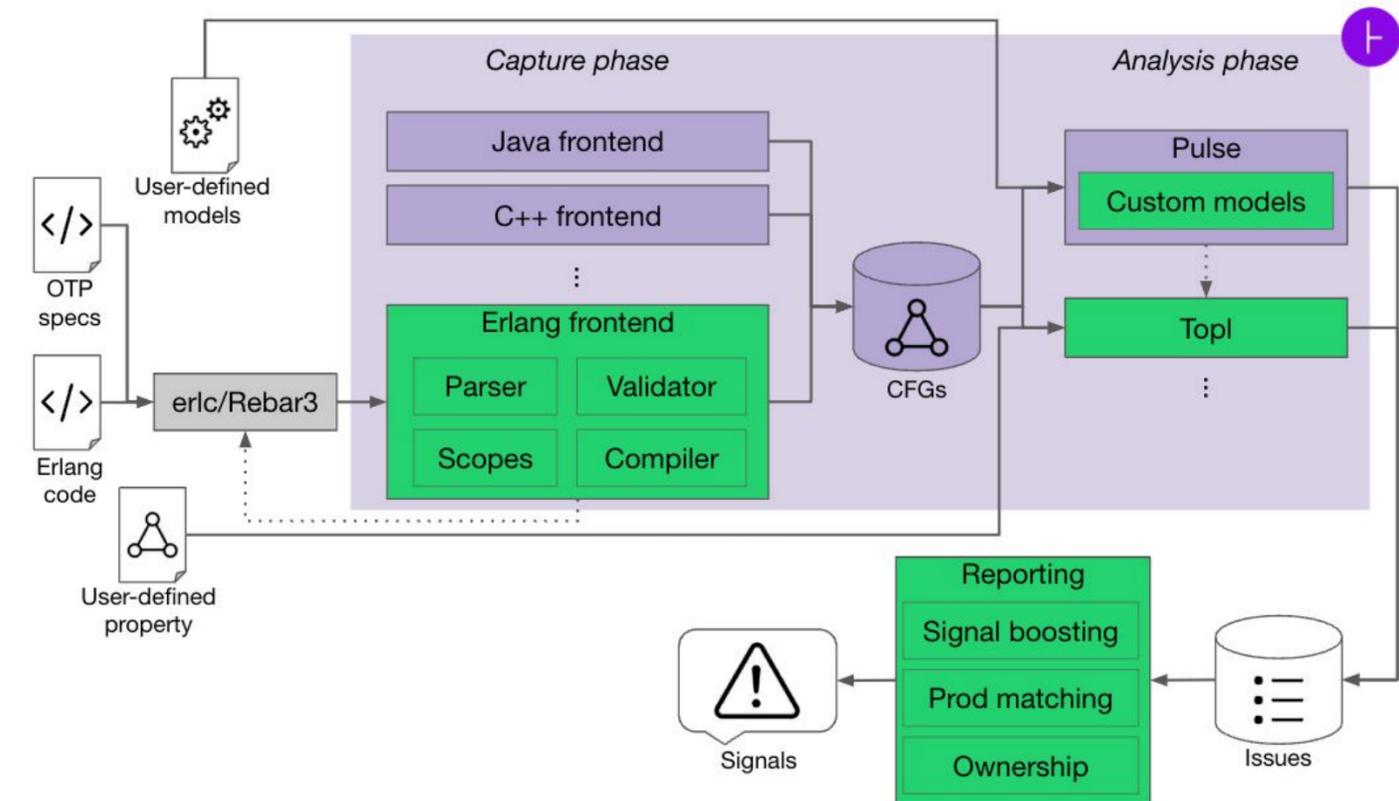
Dino Distefano
Meta
ddino@fb.com

Abstract—We introduce FAUSTA, an algorithmic traffic generation platform that enables analysis and testing at scale. FAUSTA has been deployed at Meta to analyze and test the WhatsApp plat-

into continuous integration in industry at scale (applications consisting of millions of lines of code, used by over two billion

Summary

- **InfERL**: scalable and extensible static analysis for Erlang
- **Reliability** issues + **user-defined** properties
- **Challenges**: functional, let it crash, dynamic typing, concurrent, scalability
- Erlang **frontend**
- **Extensions** to Pulse and TopI
- Promising **results** on WhatsApp server
- Available **open-source**



github.com/facebook/infer/tree/main/infer/src/erlang

